

Integrated Circuits Group

LH28F160BJHE-TTL90 Flash Memory 16M (1M × 16/2M × 8)

(Model No.: LHF16J04)

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LHF16J04

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LHF16J04

CONTENTS

PAGE

1 INTRODUCTION
1.1 Features
1.2 Product Overview
1.3 Product Description4
1.3.1 Package Pinout4
1.3.2 Block Organization4
2 PRINCIPLES OF OPERATION
2.1 Data Protection
3 BUS OPERATION
3.1 Read
3.2 Output Disable
3.3 Standby8
3.4 Reset
3.5 Read Identifier Codes9
3.6 Write
4 COMMAND DEFINITIONS9
4.1 Read Array Command12
4.2 Read Identifier Codes Command12
4.3 Read Status Register Command12
4.4 Clear Status Register Command
4.5 Block Erase Command
4.6 Full Chip Erase Command
4.7 Word/Byte Write Command
4.8 Block Erase Suspend Command14
4.9 Word/Byte Write Suspend Command14

5 DESIGN CONSIDERATIONS	25
5.1 Three-Line Output Control	25
5.2 RY/BY# and WSM Polling	25
5.3 Power Supply Decoupling	25
5.4 V _{CCW} Trace on Printed Circuit Boards	25
5.5 V _{CC} , V _{CCW} , RP# Transitions	25
5.6 Power-Up/Down Protection	
5.7 Power Dissipation	26

6 ELECTRICAL SPECIFICATIONS 27
6.1 Absolute Maximum Ratings 27
6.2 Operating Conditions
6.2.1 Capacitance
6.2.2 AC Input/Output Test Conditions
6.2.3 DC Characteristics
6.2.4 AC Characteristics - Read-Only Operations 31
6.2.5 AC Characteristics - Write Operations
6.2.6 Alternative CE#-Controlled Writes
6.2.7 Reset Operations
6.2.8 Block Erase, Full Chip Erase, Word/Byte Write and
Lock-Bit Configuration Performance 39

1

PAGE

LHF16J04



- Industry-Standard Packaging — 48-Lead TSOP
 - ETOX^{TM*} Nonvolatile Flash Technology
 - CMOS Process (P-type silicon substrate)
 - Not designed or rated as radiation hardened

SHARP's LH28F160BJHE-TTL90 Flash memory is a high-density, low-cost. nonvolatile, read/write storage solution for a wide range of applications.

LH28F160BJHE-TTL90 can operate at V_{CC}=2.7V-3.6V and V_{CCW}=2.7V-3.6V or 11.7V-12.3V. Its low voltage operation capability realize battery life and suits for cellular phone application.

Its Boot, Parameter and Main-blocked architecture, low voltage and extended cycling provide for highly flexible component suitable for portable terminals and personal computers. Its enhanced suspend capabilities provide for an ideal solution for code + data storage applications.

For secure code storage applications, such as networking, where code is either directly executed out of flash or downloaded to DRAM, the LH28F160BJHE-TTL90 offers four levels of protection: absolute protection with V_{CCW}≤V_{CCWLK}, selective hardware block locking or flexible software block locking. These alternatives give designers ultimate control of their code security needs.

The LH28F160BJHE-TTL90 is manufactured on SHARP's 0.25µm ETOX^{TM*} process technology. It come in industrystandard package: the 48-lead TSOP, ideal for board constrained applications.

*ETOX is a trademark of Intel Corporation.

Thirty-one 32K-word (64K-byte) Main Blocks

— Minimum 100,000 Block Erase Cycles

Top Boot Location

Extended Cycling Capability

1 INTRODUCTION

This datasheet contains LH28F160BJHE-TTL90 specifications. Section 1 provides a flash memory overview. Sections 2, 3, 4 and 5 describe the memory organization and functionality. Section 6 covers electrical specifications.

1.1 Features

Key enhancements of LH28F160BJHE-TTL90 boot block Flash memory are:

- •Single low voltage operation
- •Low power consumption
- •Enhanced Suspend Capabilities
- •Boot Block Architecture

Please note following:

•V_{CCWLK} has been lowered to 1.0V to support 2.7V-3.6V block erase, full chip erase, word/byte write and lock-bit configuration operations. The V_{CCW} voltage transitions to GND is recommended for designs that switch V_{CCW} off during read operation.

1.2 Product Overview

The LH28F160BJHE-TTL90 is a high-performance 16Mbit Boot Block Flash memory organized as 1M-word of 16 bits or 2M-byte of 8 bits. The 1M-word/2M-byte of data is arranged in two 4K-word/8K-byte boot blocks, six 4Kword/8K-byte parameter blocks and thirty-one 32Kword/64K-byte main blocks which are individually erasable, lockable and unlockable in-system. The memory map is shown in Figure 3.

The dedicated V_{CCW} pin gives complete data protection when $V_{CCW} \leq V_{CCWLK}$.

A Command User Interface (CUI) serves as the interface between the system processor and internal operation of the device. A valid command sequence written to the CUI initiates device automation. An internal Write State Machine (WSM) automatically executes the algorithms and timings necessary for block erase, full chip erase, word/byte write and lock-bit configuration operations. A block erase operation erases one of the device's 32K-word/64K-byte blocks typically within 1.2s (3V V_{CC} , 3V V_{CCW}). 4K-word/8K-byte blocks typically within 0.6s (3V V_{CC} , 3V V_{CCW}) independent of other blocks. Each block can be independently erased minimum 100,000 times. Block erase suspend mode allows system software to suspend block erase to read or write data from any other block.

Writing memory data is performed in word/byte increments of the device's 32K-word blocks typically within 33 μ s (3V V_{CC}. 3V V_{CCW}), 64K-byte blocks typically within 31 μ s (3V V_{CC}. 3V V_{CCW}), 4K-word blocks typically within 36 μ s (3V V_{CC}. 3V V_{CCW}), 8K-byte blocks typically within 32 μ s (3V V_{CC}. 3V V_{CCW}), 8K-byte blocks typically within 32 μ s (3V V_{CC}. 3V V_{CCW}). Word/byte write suspend mode enables the system to read data or execute code from any other flash memory array location.

Individual block locking uses a combination of bits, thirtynine block lock-bits, a permanent lock-bit and WP# pin. to lock and unlock blocks. Block lock-bits gate block erase. full chip erase and word/byte write operations, while the permanent lock-bit gates block lock-bit modification and locked block alternation. Lock-bit configuration operations (Set Block Lock-Bit, Set Permanent Lock-Bit and Clear Block Lock-Bits commands) set and cleared lock-bits.

The status register indicates when the WSM's block erase, full chip erase, word/byte write or lock-bit configuration operation is finished.

The RY/BY# output gives an additional indicator of WSM activity by providing both a hardware signal of status (versus software polling) and status masking (interrupt masking for background block erase, for example). Status polling using RY/BY# minimizes both CPU overhead and system power consumption. When low, RY/BY# indicates that the WSM is performing a block erase, full chip erase, word/byte write or lock-bit configuration. RY/BY#-high Z indicates that the WSM is ready for a new command, block erase is suspended (and word/byte write is inactive), word/byte write is suspended. or the device is in reset mode.

LHF16J04

The access time is 90ns (t_{AVQV}) over the operating temperature range (-40°C to +85°C) and V_{CC} supply voltage range of 2.7V-3.6V.

The Automatic Power Savings (APS) feature substantially reduces active current when the device is in static mode (addresses not switching). In APS mode, the typical I_{CCR} current is 2µA (CMOS) at 3.0V V_{CC} .

When CE# and RP# pins are at V_{CC} , the I_{CC} CMOS standby mode is enabled. When the RP# pin is at GND, reset mode is enabled which minimizes power consumption and provides write protection. A reset time (t_{PHQV}) is required from RP# switching high until outputs are valid. Likewise, the device has a wake time (t_{PHEL}) from RP#-high until writes to the CUI are recognized. With RP# at GND, the WSM is reset and the status register is cleared.

Please do not execute reprogramming "0" for the bit which has already been programed "0". Overwrite operation may generate unerasable bit. In case of reprogramming "0" to the data which has been programed "1".

Program "0" for the bit in which you want to change data from "1" to "0".

Program "1" for the bit which has already been programmed "0".

For example, changing data from "10111101" to "10111100" requires "11111110" programming.

1.3 Product Description

1.3.1 Package Pinout

LH28F160BJHE-TTL90 Boot Block Flash memory is available in 48-lead TSOP package (see Figure 2).

1.3.2 Block Organization

This product features an asymmetrically-blocked architecture providing system memory integration. Each erase block can be erased independently of the others up to 100,000 times. For the address locations of the blocks, see the memory map in Figure 3.

Boot Blocks: The boot block is intended to replace a dedicated boot PROM in a microprocessor or microcontroller-based system. This boot block 4K words (4.096words) features hardware controllable write-protection to protect the crucial microprocessor boot code from accidental modification. The protection of the boot block is controlled using a combination of the V_{CCW} , RP#, WP# pins and block lock-bit.

Parameter Blocks: The boot block architecture includes parameter blocks to facilitate storage of frequently update small parameters that would normally require an EEPROM. By using software techniques, the word-rewrite functionality of EEPROMs can be emulated. Each boot block component contains six parameter blocks of 4K words (4,096 words) each. The protection of the parameter block is controlled using a combination of the V_{CCW}, RP# and block lock-bit.

Main Blocks: The reminder is divided into main blocks for data or code storage. Each 16M-bit device contains thirty-one 32K words (32,768 words) blocks. The protection of the main block is controlled using a combination of the V_{CCW} , RP# and block lock-bit.



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Sumah al	Turn	Name and Function
Symbol	Туре	
A ₋₁ A ₀ -A ₁₉	INPUT	ADDRESS INPUTS: Inputs for addresses during read and write operations. Addresses are internally latched during a write cycle. A ₋₁ : Lower address input while BYTE# is V_{IL} . A ₋₁ pin changes DQ_{15} pin while BYTE# is V_{IH} . A ₁₅ -A ₁₉ : Main Block Address. A ₁₂ -A ₁₉ : Boot and Parameter Block Address.
DQ ₀ -DQ ₁₅	INPUT/ OUTPUT	DATA INPUT/OUTPUTS: Inputs data and commands during CUI write cycles: outputs data during memory array. status register and identifier code read cycles. Data pins float to high-impedance when the chip is deselected or outputs are disabled. Data is internally latched during a write cycle. DQ_8 - DQ_{15} pins are not used while byte mode (BYTE#= V_{IL}). Then, DQ_{15} pin changes A ₁ address input.
CE#	INPUT	CHIP ENABLE: Activates the device's control logic, input buffers, decoders and sense amplifier. CE#-high deselects the device and reduces power consumption to standby levels.
RP#	INPUT	RESET: Resets the device internal automation. RP#-high enables normal operation. When driven low, RP# inhibits write operations which provides data protection during power transitions. Exit from reset mode sets the device to read array mode. RP# must be V_{IL} during power-up.
OE#	INPUT	OUTPUT ENABLE: Gates the device's outputs during a read cycle.
WE#	INPUT	WRITE ENABLE: Controls writes to the CUI and array blocks. Addresses and data are latched o the rising edge of the WE# pulse.
WP#	INPUT	WRITE PROTECT: When WP# is V_{IL} , boot blocks cannot be written or erased. When WP# is V_{IH} , locked boot blocks can not be written or erased. WP# is not affected parameter and main blocks.
BYTE#	INPUT	BYTE ENABLE: BYTE# V_{IL} places device in byte mode (×8). All data is then input or output or DQ_{0-7} , and DQ_{8-15} float. BYTE# V_{IH} places the device in word mode (×16), and turns off the A_input buffer.
RY/BY#	OPEN DRAIN OUTPUT	READY/BUSY#: Indicates the status of the internal WSM. When low, the WSM is performing an internal operation (block erase, full chip erase, word/byte write or lock-bit configuration). RY/BY#-high Z indicates that the WSM is ready for new commands, block erase is suspended, and word/byte write is inactive, word/byte write is suspended, or the device is in reset mode.
V _{ccw}	SUPPLY	BLOCK ERASE, FULL CHIP ERASE, WORD/BYTE WRITE OR LOCK-BIT CONFIGURATION POWER SUPPLY: For erasing array blocks, writing words/bytes or configuring lock-bits. With $V_{CCW} \leq V_{CCWLK}$, memory contents cannot be altered. Block erase, fu chip erase, word/byte write and lock-bit configuration with an invalid V_{CCW} (see 6.2.3 DC Characteristics) produce spurious results and should not be attempted. Applying 12V±0.3V to V_{CCW} during erase/write can only be done for a maximum of 1000 cycles on each block. V_{CCW} may be connected to 12V±0.3V for a total of 80 hours maximum.
v _{cc}	SUPPLY	DEVICE POWER SUPPLY: Do not float any power pins. With $V_{CC} \leq V_{LKO}$, all write attempts to the flash memory are inhibited. Device operations at invalid V_{CC} voltage (see 6.2.3 DC Characteristics) produce spurious results and should not be attempted.
GND	SUPPLY	GROUND: Do not float any ground pins.
NC		NO CONNECT: Lead is not internal connected; it may be driven or floated.

2 PRINCIPLES OF OPERATION

The LH28F160BJHE-TTL90 flash memory includes an on-chip WSM to manage block erase. full chip erase. word/byte write and lock-bit configuration functions. It allows for: 100% TTL-level control inputs. fixed power supplies during block erase, full chip erase. word/byte write and lock-bit configuration. and minimal processor overhead with RAM-like interface timings.

After initial device power-up or return from reset mode (see section 3 Bus Operations), the device defaults to read array mode. Manipulation of external memory control pins allow array read, standby and output disable operations.

Status register and identifier codes can be accessed through the CUI independent of the V_{CCW} voltage. High voltage on V_{CCW} enables successful block erase, full chip erase, word/byte write and lock-bit configurations. All functions associated with altering memory contents-block erase, full chip erase, word/byte write, lock-bit configuration, status and identifier codes-are accessed via the CUI and verified through the status register.

Commands are written using standard microprocessor write timings. The CUI contents serve as input to the WSM, which controls the block erase. full chip erase, word/byte write and lock-bit configuration. The internal algorithms are regulated by the WSM, including pulse repetition, internal verification and margining of data. Addresses and data are internally latched during write cycles. Writing the appropriate command outputs array data, accesses the identifier codes or outputs status register data.

Interface software that initiates and polls progress of block erase, full chip erase, word/byte write and lock-bit configuration can be stored in any block. This code is copied to and executed from system RAM during flash memory updates. After successful completion, reads are again possible via the Read Array command. Block erase suspend allows system software to suspend a block erase to read/write data from/to blocks other than that which is suspend. Word/byte write suspend allows system software to suspend a word/byte write to read data from any other flash memory array location.

[A ₁₉ -A ₀]	Top Boot		[A ₁₉ -A ₋₁]
FFFFF FF000	4KW/SKB Boot Block	0	IFFFFF IFE000
FEFFF FE000	4KW/8KB Boot Block	1	IFDFFF 1FC000
FDFFF	4KW/8KB Parameter Block	0	IFBFFF
FD000 FCFFF	4KW/8KB Parameter Block	1	1FA000 1F9FFF
FC000 FBFFF	4KW/8KB Parameter Block	2	1F8000 1F7FFF
FB000 FAFFF	4KW/8KB Parameter Block	3	1F6000 1F5FFF
FA000 F9FFF	4KW/8KB Parameter Block		1F4000 1F3FFF
F9000 F8FFF	4KW/8KB Parameter Block	5	IF2000 IF1FFF
F8000 F7FFF		0	1F0000 IEFFFF
F0000 EFFFF	32KW/64KB Main Block		1E0000 1DFFFF
E8000 E7FFF	32KW/64KB Main Block	1	1D0000 ICFFFF
E0000 DFFFF	32KW/64KB Main Block	2	1C0000 1BFFFF
D8000	32KW/64KB Main Block	3	1B0000 IAFFFF
D0000 CFFFF	32KW/64KB Main Block	4	1A0000 19FFFF
C8000	32KW/64KB Main Block	5	190000
C7FFF C0000	32KW/64KB Main Block	6	18FFFF 180000
BFFFF B8000	32KW/64KB Main Block	7	17FFFF 170000
B7FFF B0000	32KW/64KB Main Block	8	16FFFF 160000
AFFFF A8000	32KW/64KB Main Block	9	15FFFF 150000
A7FFF A0000	32KW/64KB Main Block	10	14FFFF 140000
9FFFF 98000	32KW/64KB Main Block	11	13FFFF 130000
97FFF	32KW/64KB Main Block	12	12FFFF 120000
90000 8FFFF	32KW/64KB Main Block	13	IIFFFF
88000 87FFF	32KW/64KB Main Block	14	110000 10FFFF
80000 7FFFF	32KW/64KB Main Block	15	100000 0FFFFF
78000 77FFF	32KW/64KB Main Block	16	0F0000 0EFFFF
70000 6FFFF	32KW/64KB Main Block	17	0E0000 0DFFFF
68000 67FFF	32KW/64KB Main Block	18	0D0000 0CFFFF
60000 5FFFF			0C0000 0BFFFF
58000 57FFF	32KW/64KB Main Block	19	0B0000 0AFFFF
50000 4FFFF	32KW/64KB Main Block	20	0A0000 09FFFF
48000 47FFF	32KW/64KB Main Block	21	090000 08FFFF
40000 3FFFF	32KW/64KB Main Block	22	080000 07FFFF
38000	32KW/64KB Main Block	23	070000 06FFFF
37FFF 30000	32KW/64KB Main Block	24	060000 05FFFF
2FFFF 28000	32KW/64KB Main Block	25	050000
27FFF 20000	32KW/64KB Main Block	26	04FFFF 040000
1FFFF 18000	32KW/64KB Main Block	27	03FFFF 030000
17FFF 10000	32KW/64KB Main Block	28	02FFFF 020000
0FFFF 08000	32KW/64KB Main Block	29	01FFFF 010000
07FFF 00000	32KW/64KB Main Block	30	00FFFF

Figure 3. Memory Map

2.1 Data Protection

When $V_{CCW} \leq V_{CCWLK}$, memory contents cannot be altered. The CUI, with two-step block erase, full chip erase, word/byte write or lock-bit configuration command sequences, provides protection from unwanted operations even when high voltage is applied to V_{CCW} . All write functions are disabled when V_{CC} is below the write lockout voltage V_{LKO} or when RP# is at V_{IL} . The device's block locking capability provides additional protection from inadvertent code or data alteration by gating block erase, full chip erase and word/byte write operations. Refer to Table 5 for write protection alternatives.

3 BUS OPERATION

The local CPU reads and writes flash memory in-system. All bus cycles to or from the flash memory conform to standard microprocessor bus cycles.

3.1 Read

Information can be read from any block, identifier codes or status register independent of the V_{CCW} voltage. RP# can be at V_{IH} .

The first task is to write the appropriate read mode command (Read Array, Read Identifier Codes or Read Status Register) to the CUI. Upon initial device power-up or after exit from reset mode, the device automatically resets to read array mode. Six control pins dictate the data flow in and out of the component: CE#, OE#, BYTE#, WE#, RP# and WP#. CE# and OE# must be driven active to obtain data at the outputs. CE# is the device selection control, and when active enables the selected memory device. OE# is the data output (DQ₀-DQ₁₅) control and when active drives the selected memory data onto the I/O bus. BYTE# is the device I/O interface mode control. WE# must be at V_{IH}, RP# must be at V_{IH}. Figure 14, 15 illustrates read cycle.

3.2 Output Disable

With OE# at a logic-high level (V_{IH}), the device outputs are disabled. Output pins (DQ_0 - DQ_{15}) are placed in a high-impedance state.

3.3 Standby

CE# at a logic-high level (V_{IH}) places the device in standby mode which substantially reduces device power consumption. DQ_0 - DQ_{15} outputs are placed in a highimpedance state independent of OE#. If deselected during block erase, full chip erase, word/byte write or lock-bit configuration, the device continues functioning, and consuming active power until the operation completes.

3.4 Reset

RP# at V_{IL} initiates the reset mode.

In read modes, RP#-low deselects the memory, places output drivers in a high-impedance state and turns off all internal circuits. RP# must be held low for a minimum of 100ns. Time t_{PHQV} is required after return from reset mode until initial memory access outputs are valid. After this wake-up interval, normal operation is restored. The CUI is reset to read array mode and status register is set to 80H.

During block erase, full chip erase, word/byte write or lock-bit configuration modes. RP#-low will abort the operation. RY/BY# remains low until the reset operation is complete. Memory contents being altered are no longer valid; the data may be partially erased or written. Time t_{PHWL} is required after RP# goes to logic-high (V_{IH}) before another command can be written.

As with any automated device, it is important to assert RP# during system reset. When the system comes out of reset, it expects to read from the flash memory. Automated flash memories provide status information when accessed during block erase, full chip erase, word/byte write or lock-bit configuration modes. If a CPU reset occurs with no flash memory reset, proper CPU initialization may not occur because the flash memory may be providing status information instead of array data. SHARP's flash memories allow proper CPU initialization following a system reset through the use of the RP# input. In this application. RP# is controlled by the same RESET# signal that resets the system CPU.

3.5 Read Identifier Codes

The read identifier codes operation outputs the manufacturer code, device code, block lock configuration codes for each block and the permanent lock configuration code (see Figure 4). Using the manufacturer and device codes, the system CPU can automatically match the device with its proper algorithms. The block lock and permanent lock configuration codes identify locked and unlocked blocks and permanent lock-bit setting.

3.6 Write

Writing commands to the CUI enable reading of device data and identifier codes. They also control inspection and clearing of the status register. When $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$, the CUI additionally controls block erase, full chip erase, word/byte write and lock-bit configuration.

The Block Erase command requires appropriate command data and an address within the block to be erased. The Full Chip Erase command requires appropriate command data and an address within the device. The Word/Byte Write command requires the command and address of the location to be written. Set Permanent and Block Lock-Bit commands require the command and address within the device (Permanent Lock) or block within the device (Block Lock) to be locked. The Clear Block Lock-Bits command requires the command and address within the device.

The CUI does not occupy an addressable memory location. It is written when WE# and CE# are active. The address and data needed to execute a command are latched on the rising edge of WE# or CE# (whichever goes high first). Standard microprocessor write timings are used. Figures 16 and 17 illustrate WE# and CE# controlled write operations.

4 COMMAND DEFINITIONS

When the V_{CCW} voltage \leq V_{CCWLK}. Read operations from the status register, identifier codes, or blocks are enabled. Placing V_{CCWH1/2} on V_{CCW} enables successful block erase, full chip erase, word/byte write and lock-bit configuration operations.

Device operations are selected by writing specific commands into the CUI. Table 3 defines these commands.

	Top Boot
[A ₁₉ -A ₀]*	•
FFFFF	Reserved for Future Implementation
FF003	
FF002	Boot Block 0 Lock Configuration Code
FF001 FF000	Reserved for Future Implementation Boot Block 0
FEFFF	Reserved for Future Implementation
FE003	
FE002	Boot Block 1 Lock Configuration Code
FE001 FE000	Reserved for Future Implementation Boot Block 1
FDFFF	Reserved for Future Implementation
FD002	Parameter Block 0 Lock Configuration Code
FD001	Reserved for Future Implementation
FD000	Parameter Block 0
FCFFF F9000	(Parameter Blocks 1 through 4)
F8FFF	Reserved for Future Implementation
F8003 F8002	Parameter Block 5 Lock Configuration Code
F8001 F8000	Reserved for Future Implementation Parameter Block 5
F7FFF	Reserved for Future Implementation
F0003	
F0002	Main Block 0 Lock Configuration Code
F0001	Reserved for Future Implementation
F0000	Main Block 0
EFFFF 08000	(Main Blocks 1 through 29)
07FFF	Reserved for Future Implementation
00004	
00003	Permanent Lock Configuration Code
00002	Main Block 30 Lock Configuration Code
00001	Device Code
00000	Manufacturer Code Main Block 30

Figure 4. Device Identifier Code Memory Map

LHF16J04

Mode	Notes	RP#	CE#	OE#	WE#	Address	V _{CCW}	DQ ₀₋₁₅	RY/BY# ⁽³
Read	8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	X	X	D _{OUT}	X
Output Disable		V _{IH}	V _{IL}	V _{IH}	V _{IH}	X	X	High Z	X
Standby		V _{IH}	V _{IH}	X	X	X	X	High Z	X
Reset	4	V _{IL}	X	X	X	X	X	High Z	High Z
Read Identifier Codes	8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	See Figure 4	х	Note 5	High Z
Write	6,7,8	V _{IH}	VIL	V _{IH}	V _{IL}	X	Х	D _{IN}	X

Table 2.2. Bus Operations $(BYTE \#=V_{III})^{(1,2)}$

	Tuoto B.		(Cranono (
Notes	RP#	CE#	OE#	WE#	Address	V _{CCW}	DQ ₀₋₇	RY/BY# ⁽³
8	V _{IH}	VIL	V _{IL}	V _{IH}	X	X	D _{OUT}	X
					X	X	High Z	X
			X	X	X	X	High Z	·X
4		X	X	X	X	X	High Z	High Z
8	V _{IH}	V _{IL}	V _{IL}	V _{IH}	See Figure 4	x	Note 5	High Z
6,7.8	V _{IH}	V _{IL}	V _{IH}	V _{IL}	X	Х	D _{IN}	X
	8 4 8	Notes RP# 8 V _{IH} VIH V _{IH} 4 V _{IL} 8 V _{IH}	$\begin{tabular}{ c c c c c c c } \hline Notes & RP# & CE# \\ \hline 8 & V_{IH} & V_{IL} \\ \hline & V_{IH} & V_{IL} \\ \hline & V_{IH} & V_{IL} \\ \hline & 4 & V_{IL} & X \\ \hline 8 & V_{IH} & V_{IL} \\ \hline \end{tabular}$	$\begin{tabular}{ c c c c c c c c c c c } \hline Notes & RP# & CE# & OE# \\ \hline & 8 & V_{IH} & V_{IL} & V_{IL} \\ \hline & V_{IH} & V_{IL} & V_{IH} \\ \hline & V_{IH} & V_{IH} & X \\ \hline & 4 & V_{IL} & X & X \\ \hline & 8 & V_{IH} & V_{IL} & V_{IL} \\ \hline \end{array}$	$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	$\begin{tabular}{ c c c c c c c c c c c c c c c c } \hline Notes & RP# & CE# & OE# & WE# & Address \\ \hline 8 & V_{IH} & V_{IL} & V_{IL} & V_{IH} & X \\ \hline & V_{IH} & V_{IL} & V_{IH} & V_{IH} & X \\ \hline & V_{IH} & V_{IL} & V_{IH} & X & X \\ \hline 4 & V_{IL} & X & X & X \\ \hline 8 & V_{IH} & V_{IL} & V_{IL} & V_{IH} & See \\ \hline Figure 4 \end{tabular}$	$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	$ \begin{array}{ c c c c c c c c c c c c c c c c c c c$

NOTES:

Refer to DC Characteristics. When V_{CCW}≤V_{CCWLK}, memory contents can be read, but not altered.
 X can be V_{IL} or V_{IH} for control pins and addresses. and V_{CCWLK} or V_{CCWH1/2} for V_{CCW}. See DC Characteristics for

 RY/BY# is V_{OL} when the WSM is executing internal block erase, full chip erase, word/byte write or lock-bit configuration algorithms. It is High Z during when the WSM is not busy, in block erase suspend mode (with word/byte write inactive), word/byte write suspend mode or reset mode.

4. RP# at GND±0.2V ensures the lowest power consumption.

5. See Section 4.2 for read identifier code data.

6. Command writes involving block erase, full chip erase, word/byte write or lock-bit configuration are reliably executed when $V_{CCW}=V_{CCWH1/2}$ and $V_{CC}=2.7V-3.6V$. 7. Refer to Table 3 for valid D_{IN} during a write operation.

8. Never hold OE# low and WE# low at the same timing.



	Bus Cycles		Fi	First Bus Cycle			Second Bus Cycle		
Command	Req [*] d.	Notes	Oper ⁽¹⁾	Addr ⁽²⁾	Data ⁽³⁾	Oper ⁽¹⁾	Addr ⁽²⁾	Data ⁽³⁾	
Read Array/Reset	1		Write	Х	FFH				
Read Identifier Codes	≥2	4	Write	Х	90H	Read	IA	ID	
Read Status Register	2		Write	X	70 H	Read	Х	SRD	
Clear Status Register	1		Write	Х	50H				
Block Erase	2	5	Write	Х	20H	Write	BA	D0H	
Full Chip Erase	2		Write	X	30H	Write	Х	D0H	
Word/Byte Write	2	5,6	Write	Х	40H or 10H	Write	WA	WD	
Block Erase and Word/Byte Write Suspend	1	5	Write	Х	BOH				
Block Erase and Word/Byte Write Resume	1	5	Write	х	D0H				
Set Block Lock-Bit	2	8	Write	Х	60H	Write	BA	01H	
Clear Block Lock-Bits	2	7.8	Write	Х	60H	Write	X	D0H	
Set Permanent Lock-Bit	2	9	Write	Х	60H	Write	Х	FlH	

NOTES:

1. BUS operations are defined in Table 2.1 and Table 2.2.

2. X=Any valid address within the device.

IA=Identifier Code Address: see Figure 4.

BA=Address within the block being erased.

WA=Address of memory location to be written.

3. SRD=Data read from status register. See Table 6 for a description of the status register bits.

WD=Data to be written at location WA. Data is latched on the rising edge of WE# or CE# (whichever goes high first). ID=Data read from identifier codes.

4. Following the Read Identifier Codes command, read operations access manufacturer, device, block lock configuration and permanent lock configuration codes. See Section 4.2 for read identifier code data.

5. If WP# is V_{IL}, boot blocks are locked without block lock-bits state. If WP# is V_{IH}, boot blocks are locked by block lock-bits. The parameter and main blocks are locked by block lock-bits without WP# state.

6. Either 40H or 10H are recognized by the WSM as the word/byte write setup.

7. The clear block lock-bits operation simultaneously clears all block lock-bits.

8. If the permanent lock-bit is set. Set Block Lock-Bit and Clear Block Lock-Bits commands can not be done.

9. Once the permanent lock-bit is set, permanent lock-bit reset is unable.

10. Commands other than those shown above are reserved by SHARP for future device implementations and should not be used.

4.1 Read Array Command

Upon initial device power-up and after exit from reset mode, the device defaults to read array mode. This operation is also initiated by writing the Read Array command. The device remains enabled for reads until another command is written. Once the internal WSM has started a block erase, full chip erase. word/byte write or lock-bit configuration the device will not recognize the Read Array command until the WSM completes its operation unless the WSM is suspended via an Erase Suspend or Word/Byte Write Suspend command. The Read Array command functions independently of the V_{CCW} voltage and RP# can be V_{IH} .

4.2 Read Identifier Codes Command

The identifier code operation is initiated by writing the Read Identifier Codes command. Following the command write, read cycles from addresses shown in Figure 4 retrieve the manufacturer, device, block lock configuration and permanent lock configuration codes (see Table 4 for identifier code values). To terminate the operation, write another valid command. Like the Read Array command, the Read Identifier Codes command functions independently of the V_{CCW} voltage and RP# can be V_{IH}. Following the Read Identifier Codes command, the following information can be read:

Table 4. Identifier Codes									
Code	Address ⁽²⁾	Data ⁽³⁾							
	$[A_{19} - A_0]$	$[DQ_7 - DQ_0]$							
Manufacture Code	00000H	B0H							
Device Code	00001H	E8H							
Block Lock Configuration	BA ⁽¹⁾ +2								
•Block is Unlocked		DQ ₀ =0							
•Block is Locked		DQ ₀ =1							
•Reserved for Future Use		DQ ₁₋₇							
Permanent Lock Configuration	00003H								
•Device is Unlocked		DQ ₀ =0							
•Device is Locked		$DQ_0 = 1$							
•Reserved for Future Use		DQ ₁₋₇							

Table 4. Identifier Codes

NOTE:

1. BA selects the specific block lock configuration code to be read. See Figure 4 for the device identifier code memory map.

2. A₋₁ don't care in byte mode.

3. $D\dot{Q}_{15}$ - DQ_8 outputs 00H in word mode.

4.3 Read Status Register Command

The status register may be read to determine when a block erase, full chip erase. word/byte write or lock-bit configuration is complete and whether the operation completed successfully. It may be read at any time by writing the Read Status Register command. After writing this command, all subsequent read operations output data from the status register until another valid command is written. The status register contents are latched on the falling edge of OE# or CE#, whichever occurs. OE# or CE# must toggle to V_{IH} before further reads to update the status register latch. The Read Status Register command functions independently of the V_{CCW} voltage. RP# can be V_{IH}.

4.4 Clear Status Register Command

Status register bits SR.5. SR.4. SR.3 or SR.1 are set to "1"s by the WSM and can only be reset by the Clear Status Register command. These bits indicate various failure conditions (see Table 6). By allowing system software to reset these bits, several operations (such as cumulatively erasing multiple blocks or writing several words/bytes in sequence) may be performed. The status register may be polled to determine if an error occurred during the sequence.

To clear the status register, the Clear Status Register command (50H) is written. It functions independently of the applied V_{CCW} Voltage. RP# can be V_{IH} . This command is not functional during block erase or word/byte write suspend modes.

4.5 Block Erase Command

Erase is executed one block at a time and initiated by a two-cycle command. A block erase setup is first written, followed by an block erase confirm. This command sequence requires appropriate sequencing and an address within the block to be erased (erase changes all block data to FFFFH/FFH). Block preconditioning, erase, and verify are handled internally by the WSM (invisible to the system). After the two-cycle block erase sequence is written, the device automatically outputs status register data when read (see Figure 5). The CPU can detect block erase completion by analyzing the output data of the RY/BY# pin or status register bit SR.7.

When the block erase is complete, status register bit SR.5 should be checked. If a block erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read status register mode until a new command is issued.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Block Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable block erasure can only occur when $V_{CC}=2.7V$ -3.6V and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, block contents are protected against erasure. If block erase is attempted while $V_{CCW} \leq V_{CCWLK}$. SR.3 and SR.5 will be set to "1". Successful block erase requires for boot blocks that WP# is V_{IH} and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If block erase is attempted when the excepting above conditions, SR.1 and SR.5 will be set to "1".

4.6 Full Chip Erase Command

This command followed by a confirm command erases all of the unlocked blocks. A full chip erase setup (30H) is first written, followed by a full chip erase confirm (D0H). After a confirm command is written, device erases the all unlocked blocks block by block. This command sequence requires appropriate sequencing. Block preconditioning, erase and verify are handled internally by the WSM (invisible to the system). After the two-cycle full chip erase sequence is written, the device automatically outputs status register data when read (see Figure 6). The CPU can detect full chip erase completion by analyzing the output data of the RY/BY# pin or status register bit SR.7.

When the full chip erase is complete, status register bit SR.5 should be checked. If erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read

status register mode until a new command is issued. If error is detected on a block during full chip erase operation, WSM stops erasing. Full chip erase operation start from lower address block, finish the higher address block. Full chip erase can not be suspended.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Full Chip Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable full chip erasure can only occur when V_{CC} =2.7V-3.6V and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, block contents are protected against erasure. If full chip erase is attempted while $V_{CCW} \leq V_{CCWLK}$, SR.3 and SR.5 will be set to "1". Successful full chip erase requires for boot blocks that WP# is V_{IH} and the corresponding block lock-bit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If all blocks are locked, SR.1 and SR.5 will be set to "1".

4.7 Word/Byte Write Command

Word/Byte write is executed by a two-cycle command sequence. Word/Byte write setup (standard 40H or alternate 10H) is written, followed by a second write that specifies the address and data (latched on the rising edge of WE#). The WSM then takes over, controlling the word/byte write and write verify algorithms internally. After the word/byte write sequence is written, the device automatically outputs status register data when read (see Figure 7). The CPU can detect the completion of the word/byte write event by analyzing the RY/BY# pin or status register bit SR.7.

When word/byte write is complete, status register bit SR.4 should be checked. If word/byte write error is detected, the status register should be cleared. The internal WSM verify only detects errors for "1"s that do not successfully write to "0"s. The CUI remains in read status register mode until it receives another command.

Reliable word/byte writes can only occur when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, memory contents are protected against word/byte writes. If word/byte write is attempted while $V_{CCW}\leq V_{CCWLK}$, status register bits SR.3 and SR.4 will be set to "1". Successful word/byte write requires for boot blocks that WP# is V_{IH} and the corresponding block lockbit be cleared. In parameter and main blocks case, it must be cleared the corresponding block lock-bit. If word/byte write is attempted when the excepting above conditions, SR.1 and SR.4 will be set to "1".

4.8 Block Erase Suspend Command

The Block Erase Suspend command allows block-erase interruption to read or word/byte write data in another block of memory. Once the block erase process starts, writing the Block Erase Suspend command requests that the WSM suspend the block erase sequence at a predetermined point in the algorithm. The device outputs status register data when read after the Block Erase Suspend command is written. Polling status register bits SR.7 and SR.6 can determine when the block erase operation has been suspended (both will be set to "1"). RY/BY# will also transition to High Z. Specification t_{WHRZ2} defines the block erase suspend latency.

When Block Erase Suspend command write to the CUI, if block erase was finished, the device places read array mode. Therefore, after Block Erase Suspend command write to the CUI, Read Status Register command (70H) has to write to CUI, then status register bit SR.6 should be checked for places the device in suspend mode.

At this point, a Read Array command can be written to read data from blocks other than that which is suspended. A Word/Byte Write command sequence can also be issued during erase suspend to program data in other blocks. Using the Word/Byte Write Suspend command (see Section 4.9), a word/byte write operation can also be suspended. During a word/byte write operation with block erase suspended, status register bit SR.7 will return to "0" and the RY/BY# output will transition to V_{OL} . However, SR.6 will remain "1" to indicate block erase suspend status.

The only other valid commands while block erase is suspended are Read Status Register and Block Erase Resume. After a Block Erase Resume command is written to the flash memory, the WSM will continue the block erase process. Status register bits SR.6 and SR.7 will automatically clear and RY/BY# will return to V_{OL} . After the Erase Resume command is written, the device automatically outputs status register data when read (see Figure 8). V_{CCW} must remain at $V_{CCWH1/2}$ (the same V_{CCW} level used for block erase) while block erase is suspended. RP# must also remain at V_{IH} . WP# must also remain at V_{IL} or V_{IH} (the same WP# level used for block erase). Block erase cannot resume until word/byte write operations initiated during block erase suspend have completed.

If the period of from Block Erase Resume command write to the CUI till Block Erase Suspend command write to the CUI be short and done again and again, erase time be prolonged.

4.9 Word/Byte Write Suspend Command

The Word/Byte Write Suspend command allows word/byte write interruption to read data in other flash memory locations. Once the word/byte write process starts, writing the Word/Byte Write Suspend command requests that the WSM suspend the Word/Byte write sequence at a predetermined point in the algorithm. The device continues to output status register data when read after the Word/Byte Write Suspend command is written. Polling status register bits SR.7 and SR.2 can determine when the word/byte write operation has been suspended (both will be set to "1"). RY/BY# will also transition to High Z. Specification t_{WHRZ1} defines the word/byte write suspend latency.

When Word/Byte Write Suspend command write to the CUI, if word/byte write was finished, the device places read array mode. Therefore, after Word/Byte Write Suspend command write to the CUI, Read Status Register command (70H) has to write to CUI, then status register bit SR.2 should be checked for places the device in suspend mode.

At this point, a Read Array command can be written to read data from locations other than that which is suspended. The only other valid commands while word/byte write is suspended are Read Status Register and Word/Byte Write Resume. After Word/Byte Write Resume command is written to the flash memory, the WSM will continue the word/byte write process. Status register bits SR.2 and SR.7 will automatically clear and RY/BY# will return to V_{OL} . After the Word/Byte Write Resume command is written, the device automatically outputs status register data when read (see Figure 9). V_{CCW} must remain at $V_{CCWH1/2}$ (the same V_{CCW} level used for word/byte write) while in word/byte write suspend mode. RP# must also remain at V_{IH} . WP# must also remain at V_{IL} or V_{IH} (the same WP# level used for word/byte write).

If the period of from Word/Byte Write Resume command write to the CUI till Word/Byte Write Suspend command write to the CUI be short and done again and again, write time be prolonged.

4.10 Set Block and Permanent Lock-Bit Commands

A flexible block locking and unlocking scheme is enabled via a combination of block lock-bits. a permanent lock-bit and WP# pin. The block lock-bits and WP# pin gates program and erase operations while the permanent lock-bit gates block-lock bit modification. With the permanent lock-bit not set, individual block lock-bits can be set using the Set Block Lock-Bit command. The Set Permanent Lock-Bit command, sets the permanent lock-bit. After the permanent lock-bit is set, block lock-bits and locked block contents cannot altered. See Table 5 for a summary of hardware and software write protection options.

Set block lock-bit and permanent lock-bit are executed by a two-cycle command sequence. The set block or permanent lock-bit setup along with appropriate block or device address is written followed by either the set block lock-bit confirm (and an address within the block to be locked) or the set permanent lock-bit confirm (and any device address). The WSM then controls the set lock-bit algorithm. After the sequence is written, the device automatically outputs status register data when read (see Figure 10). The CPU can detect the completion of the set lock-bit event by analyzing the RY/BY# pin output or status register bit SR.7.

When the set lock-bit operation is complete, status register bit SR.4 should be checked. If an error is detected, the status register should be cleared. The CUI will remain in read status register mode until a new command is issued.

This two-step sequence of set-up followed by execution ensures that lock-bits are not accidentally set. An invalid Set Block or Permanent Lock-Bit command will result in status register bits SR.4 and SR.5 being set to "1". Also, reliable operations occur only when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. In the absence of this high voltage, lock-bit contents are protected against alteration.

A successful set block lock-bit operation requires that the permanent lock-bit be cleared. If it is attempted with the permanent lock-bit set. SR.1 and SR.4 will be set to "1" and the operation will fail.

4.11 Clear Block Lock-Bits Command

All set block lock-bits are cleared in parallel via the Clear Block Lock-Bits command. With the permanent lock-bit not set, block lock-bits can be cleared using only the Clear Block Lock-Bits command. If the permanent lock-bit is set, block lock-bits cannot cleared. See Table 5 for a summary of hardware and software write protection options.

Clear block lock-bits operation is executed by a two-cycle command sequence. A clear block lock-bits setup is first written. After the command is written, the device automatically outputs status register data when read (see Figure 11). The CPU can detect completion of the clear block lock-bits event by analyzing the RY/BY# Pin output or status register bit SR.7.

When the operation is complete, status register bit SR.5 should be checked. If a clear block lock-bit error is detected, the status register should be cleared. The CUI will remain in read status register mode until another command is issued.

This two-step sequence of set-up followed by execution ensures that block lock-bits are not accidentally cleared. An invalid Clear Block Lock-Bits command sequence will result in status register bits SR.4 and SR.5 being set to "1". Also, a reliable clear block lock-bits operation can only occur when $V_{CC}=2.7V-3.6V$ and $V_{CCW}=V_{CCWH1/2}$. If a clear block lock-bits operation is attempted while $V_{CCW}\leq V_{CCWLK}$, SR.3 and SR.5 will be set to "1". In the absence of this high voltage, the block lock-bits content are protected against alteration. A successful clear block lock-bits operation requires that the permanent lock-bit is not set. If it is attempted with the permanent lock-bit set, SR.1 and SR.5 will be set to "1" and the operation will fail.

If a clear block lock-bits operation is aborted due to V_{CCW} or V_{CC} transitioning out of valid range or RP# active transition, block lock-bit values are left in an undetermined state. A repeat of clear block lock-bits is required to initialize block lock-bit contents to known values. Once the permanent lock-bit is set, it cannot be cleared.

4.12 Block Locking by the WP#

This Boot Block Flash memory architecture features two hardware-lockable boot blocks so that the kernel code for the system can be kept secure while other blocks are programmed or erased as necessary.

The lockable blocks are locked when $WP\#=V_{IL}$; any program or erase operation to a locked block will result in

an error, which will be reflected in the status register. For top configuration, the top two boot blocks are lockable. For the bottom configuration, the bottom two boot blocks are lockable. If WP# is V_{IH} and block lock-bit is not set, boot block can be programmed or erased normally (Unless V_{CCW} is below V_{CCWLK}). WP# is valid only two boot blocks, other blocks are not affected.

[1		r		e Prote	ction Alternatives
Operation	V _{CCW}	RP#	Permanent Lock-Bit	Block Lock-bit	WP#	Effect
Block Erase	≤V _{CCWLK}	Х	X	X	x	All Blocks Locked.
or	>V _{CCWLK}	V _{IL}	X	X	X	All Blocks Locked.
Word/Byte		VIH	X	0	V _{IL}	2 Boot Blocks Locked.
Write					V _{IH}	Block Erase and Word/Byte Write Enabled.
				1	V _{IL}	Block Erase and Word/Byte Write Disabled.
					V _{IH}	Block Erase and Word/Byte Write Disabled.
Full Chip	≤V _{CCWLK}	Х	X	Х	Х	All Blocks Locked.
Erase	>V _{CCWLK}	V _{IL}	Х	Х	X	All Blocks Locked.
		V _{IH}	Х	х	V _{IL}	All Unlocked Blocks are Erased. 2 Boot Blocks and Locked Blocks are NOT Erased.
					V _{IH}	All Unlocked Blocks are Erased. Locked Blocks are NOT Erased.
Set Block	≤V _{CCWLK}	Х	X	x	Х	Set Block Lock-Bit Disabled.
Lock-Bit	>V _{CCWLK}	VIL	Х	X	Х	Set Block Lock-Bit Disabled.
		VIH	0	Х	X	Set Block Lock-Bit Enabled.
			1	Х	Х	Set Block Lock-Bit Disabled.
Clear Block	≤V _{CCWLK}	Х	X	Х	Х	Clear Block Lock-Bits Disabled.
Lock-Bits	>V _{CCWLK}	V _{IL}	Х	Х	Х	Clear Block Lock-Bits Disabled.
		V _{IH}	0	Х	X	Clear Block Lock-Bits Enabled.
			1	Х	Х	Clear Block Lock-Bits Disabled.
Set	≤V _{CCWLK}	Х	Х	Х	Х	Set Permanent Lock-Bit Disabled.
Permanent	>V _{CCWLK}	V _{IL}	Х	X	Х	Set Permanent Lock-Bit Disabled.
Lock-Bit		V _{IH}	X	X	Х	Set Permanent Lock-Bit Enabled.

Table 5. Write Protection Alternatives

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LHF16J04

WSMS	BESS	ECBLBS	WBWSLBS	egister Definitio	WBWSS	DPS	R
7	6	5	4	3	2	1	0
					NOTES	5:	<u>.</u>
SR.7 = WRIT 1 = Ready 0 = Busy	E STATE MAC	THINE STATUS	(WSMS)	erase, word/by	# or SR.7 to deta te write or lock- alid while SR.7=	bit configuration	
1 = Block	K ERASE SUS Erase Suspende Erase in Progres		G (BESS)				
STAT 1 = Error in Lock-I 0 = Succes	US (ECBLBS) n Block Erase, H Bits	BLOCK LOCK	or Clear Block	erase or lock-b	nd SR.4 are "1"s it configuration ience was entere	attempt, an im	rase, full chip proper
 SR.4 = WORD/BYTE WRITE AND SET LOCK-BIT STATUS (WBWSLBS) 1 = Error in Word/Byte Write or Set Block/Permanent Lock-Bit 0 = Successful Word/Byte Write or Set Block/Permanent Lock-Bit 				level. The WS only after Bloc Lock-Bit Conf guaranteed to r	provide a contin M interrogates a k Erase, Full Ch iguration comma reports accurate	nd indicates the hip Erase, Word and sequences.	e V _{CCW} level I/Byte Write o SR.3 is not
$1 = V_{CCW}$ $0 = V_{CCW}$ SR.2 = WORE (WBW) 1 = Word/F 0 = Word/F SR.1 = DEVIC 1 = Block I Lock E 0 = Unlock	D/BYTE WRITE SS) Byte Write Susp Byte Write in Pr E PROTECT S Lock-Bit, Perma Detected, Operat	eration Abort E SUSPEND ST ended ogress/Complete TATUS (DPS) anent Lock-Bit a ion Abort	ed nd/or WP#	and block lock the permanent Block Erase, F Configuration depending on t set, permanent the block lock writing the Rea permanent and	provide a contin -bit and WP# va lock-bit, block k 'ull Chip Erase. ' command seque he attempted op lock-bit is set ar and permanent l ad Identifier Coc block lock-bit s	lues. The WSN ock-bit and WI Word/Byte Wri nces. It information, if the b nd/or WP# is V ock configuration les command in tatus.	A interrogates P# only after ite or Lock-Bit is the system, block lock-bit i IL. Reading ion codes after ndicates
SR.0 = RESER	VED FOR FUT	FURE ENHANC	LEMENTS (R)		d for future use he status register		masked out

LHF16J04





Bus Operation	Command	Comments
Wnte	Read Status Register	Dala=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Full Chip Erase Setup	Data=30H Addr=X
Write	Full Chip Erase Confirm	Data=D0H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Write FFH after the last operation to place device in read array mode.

LHF16J04

or	eration	Command	Comments
	andby		Check SR.3
3	andby		1=V _{CCW} Error Detect
			Check SR.1
S	andby		1=Device Protect Detect
			(All Blocks are locked)
			Check SR.4.5
5	andby		Both 1=Command Sequence Error
			Check SR.5
51	andby		1=Full Chip Erase Error
wh	ere multiple blo	ocks are erased before t	by the Clear Status Register Command in cas full status is checked. before attempting retry or other error recover
•			· · · · · · · · · · · · · · · · · · ·

LHF16J04



LHF16J04





LHF16J04

Figure 9. Word/Byte Write Suspend/Resume Flowchart

Start

V

Write 70H

Read Status

Register

SR.7=

V.

Write 60H

v

Write 01H/F1H.

Block/Device Address

Read

Status Register

SR.7=

Full Status

Check if Desired

Y. Set Lock-Bat Complete

0

0



FULL STATUS CHECK PROCEDURE



Bus Operation	Command	Comments
Standby		Check SR.3 I=V _{CCW} Error Detect
Standby		Check SR.1 1=Device Protect Detect Permanent Lock-Bit is Set (Set Block Lock-Bit Operation)
Standby		Check SR.4.5 Both 1=Command Sequence Error
Standby		Check SR.4 1=Set Lock-Bit Error

where multiple lock-bits are set before full status is checked. If error is detected, clear the Status Register before attempting retry or other error recovery.

Figure 10. Set Block and Permanent Lock-Bit Flowchart

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LHF16J04



5 DESIGN CONSIDERATIONS

5.1 Three-Line Output Control

The device will often be used in large memory arrays. SHARP provides three control inputs to accommodate multiple memory connections. Three-line control provides for:

- a. Lowest possible memory power dissipation.
- b. Complete assurance that data bus contention will not occur.

To use these control inputs efficiently, an address decoder should enable CE# while OE# should be connected to all memory devices and the system's READ# control line. This assures that only selected memory devices have active outputs while deselected memory devices are in standby mode. RP# should be connected to the system POWERGOOD signal to prevent unintended writes during system power transitions. POWERGOOD should also toggle during system reset.

5.2 RY/BY# and WSM Polling

RY/BY# is an open drain output that should be connected to V_{CC} by a pull up resistor to provides a hardware method of detecting block erase, full chip erase, word/byte write and lock-bit configuration completion. It transitions low after block erase, full chip erase, word/byte write or lock-bit configuration commands and returns to V_{OH} (while RY/BY# is pull up) when the WSM has finished executing the internal algorithm.

RY/BY# can be connected to an interrupt input of the system CPU or controller. It is active at all times. RY/BY# is also High Z when the device is in block erase suspend (with word/byte write inactive), word/byte write suspend or reset modes.

5.3 Power Supply Decoupling

Flash memory power switching characteristics require careful device decoupling. System designers are interested in three supply current issues; standby current levels, active current levels and transient peaks produced by falling and rising edges of CE# and OE#. Transient current magnitudes depend on the device outputs' capacitive and inductive loading. Two-line control and proper decoupling capacitor selection will suppress transient voltage peaks. Each device should have a 0.1µF ceramic capacitor connected between its $V_{\mbox{\scriptsize CC}}$ and GND and between its V_{CCW} and GND. These high-frequency, low inductance capacitors should be placed as close as possible to package leads. Additionally, for every eight devices, a 4.7µF electrolytic capacitor should be placed at the array's power supply connection between V_{CC} and GND. The bulk capacitor will overcome voltage slumps caused by PC board trace inductance.

5.4 V_{CCW} Trace on Printed Circuit Boards

Updating flash memories that reside in the target system requires that the printed circuit board designer pay attention to the V_{CCW} Power supply trace. The V_{CCW} pin supplies the memory cell current for word/byte writing and block erasing. Use similar trace widths and layout considerations given to the V_{CC} power bus. Adequate V_{CCW} supply traces and decoupling will decrease V_{CCW} voltage spikes and overshoots.

5.5 V_{CC}, V_{CCW}, RP# Transitions

Block erase, full chip erase, word/byte write and lock-bit configuration are not guaranteed if V_{CCW} falls outside of a valid $V_{CCWH1/2}$ range, V_{CC} falls outside of a valid 2.7V-3.6V range. or RP# $\neq V_{IH}$. If V_{CCW} error is detected, status register bit SR.3 is set to "1" along with SR.4 or SR.5, depending on the attempted operation. If RP# transitions to V_{IL} during block erase, full chip erase, word/byte write or lock-bit configuration, RY/BY# will remain low until the reset operation is complete. Then, the operation will abort and the device will enter reset mode. The aborted operation may leave data partially altered. Therefore, the command sequence must be repeated after normal operation is restored. Device power-off or RP# transitions to V_{IL} clear the status register.

The CUI latches commands issued by system software and is not altered by V_{CCW} or CE# transitions or WSM actions. Its state is read array mode upon power-up. after exit from reset mode or after V_{CC} transitions below V_{LKO} .

5.6 Power-Up/Down Protection

SHARP

The device is designed to offer protection against accidental block erase, full chip erase, word/byte write or lock-bit configuration during power transitions. Upon power-up, the device is indifferent as to which power supply (V_{CCW} or V_{CC}) powers-up first. Internal circuitry resets the CUI to read array mode at power-up.

A system designer must guard against spurious writes for V_{CC} voltages above V_{LKO} when V_{CCW} is active. Since both WE# and CE# must be low for a command write, driving either to V_{IH} will inhibit writes. The CUI's two-step command sequence architecture provides added level of protection against data alteration.

In-system block lock and unlock capability prevents inadvertent data alteration. The device is disabled while $RP\#=V_{II}$ regardless of its control inputs state.

5.7 Power Dissipation

When designing portable systems, designers must consider battery power consumption not only during device operation, but also for data retention during system idle time. Flash memory's nonvolatility increases usable battery life because data is retained when system power is removed.

5.8 Data Protection Method

Noises having a level exceeding the limit specified in the specification may be generated under specific operating conditions on some systems. Such noises, when induced onto WE# signal or power supply. may be interpreted as false commands, causing undesired memory updating. To protect the data stored in the flash memory against unwanted overwriting, systems operating with the flash memory should have the following write protect designs. as appropriate:

1) Protecting data in specific block

When a lock bit is set, the corresponding block (includes the 2 boot blocks) is protected against overwriting. By setting a WP# to low, only the 2 boot blocks can be protected against overwriting. By using this feature, the flash memory space can be divided into the program section (locked section) and data section (unlocked section). The permanent lock bit can be used to prevent false block bit setting. For further information on setting/resetting lock-bit, refer to the specification. (See chapter 4.10 and 4.11.)

2) Data protection through V_{CCW}

When the level of V_{CCW} is lower than V_{CCWLK} (lockout voltage), write operation on the flash memory is disabled. All blocks are locked and the data in the blocks are completely write protected. For the lockout voltage, refer to the specification. (See chapter 6.2.3.)

3) Data protection through RP#

When the RP# is kept low during read mode, the flash memory will be reset mode, then write protecting all blocks. When the RP# is kept low during power up and power down sequence such as voltage transition, write operation on the flash memory is disabled, write protecting all blocks. For the details of RP# control, refer to the specification. (See chapter 5.6 and 6.2.7.)

6 ELECTRICAL SPECIFICATIONS

6.1 Absolute Maximum Ratings*

Operating Temperature During Read, Block Erase. Full Chip Erase, Word/Byte Write and Lock-Bit Configuration.....-40°C to +85°C⁽¹⁾

Storage Temperature	
During under Bias	
During non Bias	65°C to +125°C

Voltage On Any Pin	
(except V_{CC} and V_{CCW})	-0.5V to V_{CC} +0.5V ⁽²⁾

 V_{CC} Supply Voltage.....-0.2V to +4.6V⁽²⁾ V_{CCW} Supply Voltage.....-0.2V to +13.0V^(2,3)

*WARNING: Stressing the device beyond the "Absolute Maximum Ratings" may cause permanent damage. These are stress ratings only. Operation beyond the "Operating Conditions" is not recommended and extended exposure beyond the "Operating Conditions" may affect device reliability.

NOTES:

- 1. Operating temperature is for extended temperature product defined by this specification.
- 2. All specified voltages are with respect to GND. Minimum DC voltage is -0.5V on input/output pins and -0.2V on V_{CC} and V_{CCW} pins. During transitions, this level may undershoot to -2.0V for periods <20ns. Maximum DC voltage on input/output pins are V_{CC} +0.5V which, during transitions, may overshoot to V_{CC} +2.0V for periods <20ns.
- Maximum DC voltage on V_{CCW} may overshoot to +13.0V for periods <20ns. Applying 12V±0.3V to V_{CCW} during erase/write can only be done for a maximum of 1000 cycles on each block. V_{CCW} may be connected to 12V±0.3V for a total of 80 hours maximum.
- 4. Output shorted for no more than one second. No more than one output shorted at a time.

6.2 Operating Conditions

	Temperature and V _{CC} Operating Conditions						
Symbol	Parameter	Min.	Max.	Unit	Test Condition		
T _A	Operating Temperature	-40	+85	°C	Ambient Temperature		
V _{CC}	V _{CC} Supply Voltage (2.7V-3.6V)	2.7	3.6	v			

6.2.1 CAPACITANCE⁽¹⁾

		T _A =+25°C,	f=1MHz		
Symbol	Parameter	Тур.	Max.	Unit	Condition
C _{IN}	Input Capacitance	7	10	pF	V _{IN} =0.0V
C _{OUT}	Output Capacitance	9	12	pF	V _{OUT} =0.0V

NOTE:

1. Sampled, not 100% tested.



6.2.3 DC CHARACTERISTICS

		DC Ch	aracteristics			
			V _{CC} =2.	7V-3.6V		Test
Sym.	Parameter	Notes	Тур.	Max.	Unit	Conditions
I _{LI}	Input Load Current	1		±0.5	μА	V _{CC} =V _{CC} Max. V _{IN} =V _{CC} or GND
I _{LO}	Output Leakage Current	1		±0.5	μΑ	V _{CC} =V _{CC} Max. V _{OUT} =V _{CC} or GND
I _{CCS}	V _{CC} Standby Current	1,3,6	2	15	μA	CMOS Level Inputs V _{CC} =V _{CC} Max. CE#=RP#=V _{CC} ±0.2V
			0.2	2	mA	TTL Level Inputs V _{CC} =V _{CC} Max. CE#=RP#=V _{IH}
I _{CCAS}	V _{CC} Auto Power-Save Current	1,5,6	2	15	μА	CMOS Level Inputs V _{CC} =V _{CC} Max. CE#=GND±0.2V
I _{CCD}	V _{CC} Reset Power-Down Current	1	2	15	μΑ	RP#=GND±0.2V I _{OUT} (RY/BY#)=0mA
I _{CCR}	V _{CC} Read Current	1,6	15	25	mA	CMOS Level Inputs V _{CC} =V _{CC} Max., CE#=GNI f=5MHz, I _{OUT} =0mA
				30	mA	TTL Level Inputs V _{CC} =V _{CC} Max., CE#=GNI f=5MHz, I _{OUT} =0mA
I _{CCW}	V _{CC} Word/Byte Write or Set Lock-	1,7	5	17	mA	V _{CCW} =2.7V-3.6V
	Bit Current		5	12	mA	V _{CCW} =11.7V-12.3V
I _{CCE}	V _{CC} Block Erase, Full Chip Erase or	1,7	4	17	mA	V _{CCW} =2.7V-3.6V
	Clear Block Lock-Bits Current		4	12	mA	V _{CCW} =11.7V-12.3V
I _{CCWS}	V _{CC} Word/Byte Write or Block Erase Suspend Current	1,2	1	6	mA	CE#=V _{IH}
I _{CCWS}	V _{CCW} Standby or Read Current	1	±2	±15	μА	V _{CCW} ≤V _{CC}
I _{CCWR}			10	200	μΑ	V _{CCW} >V _{CC}
I _{CCWAS}	V _{CCW} Auto Power-Save Current	1,5,6	0.1	5	μΑ	CMOS Level Inputs V _{CC} =V _{CC} Max. CE#=GND±0.2V
CCWD	V _{CCW} Reset Power-Down Current	1	0.1	5	μA	RP#=GND±0.2V
I _{CCWW}	V _{CCW} Word/Byte Write or Set Lock- Bit Current	1,7	12	40 30	mA mA	V _{CCW} =2.7V-3.6V V _{CCW} =11.7V-12.3V
CCWE	V _{CCW} Block Erase, Full Chip Erase	1,7	8	25	mA	V _{CCW} =2.7V-3.6V
CLWE	or Clear Block Lock-Bits Current			20	mA	V _{CCW} =11.7V-12.3V
I _{CCWWS}	V _{CCW} Word/Byte Write or Block Erase Suspend Current	1	10	200	μΑ	V _{CCW} =V _{CCWH1/2}

		T T	V _{CC} =2.	7V-3.6V		
Sym.	Parameter	Notes	Min.	Max.	Unit	Test Conditions
V _{IL}	Input Low Voltage	7	-0.5	0.8	V	
V _{IH}	Input High Voltage	7	2.0	V _{CC} +0.5	v	
V _{OL}	Output Low Voltage	3,7		0.4	v	V _{CC} =V _{CC} Min. I _{OL} =2.0mA
V _{OH1}	Output High Voltage (TTL)	7	2.4		v	V _{CC} =V _{CC} Min. I _{OH} =-2.0mA
V _{OH2}	Output High Voltage (CMOS)	7	0.85 V _{CC}		v	V _{CC} =V _{CC} Min. I _{OH} =-2.5mA
			V _{CC} -0.4		v	$V_{CC} = V_{CC} Min.$ $I_{OH} = -100 \mu A$
V _{CCWLK}	V _{CCW} Lockout during Normal Operations	4,7		1.0	v	
V _{CCWH1}	V _{CCW} during Block Erase, Full Chip Erase, Word/Byte Write or Lock-Bit		2.7	3.6	v	

V_{LKO} NOTES:

V_{CCWH2}

8

1. All currents are in RMS unless otherwise noted. Typical values at nominal V_{CC} voltage and $T_A = +25^{\circ}C$. 2. I_{CCWS} and I_{CCES} are specified with the device de-selected. If read or word/byte written while in erase suspend mode, the

11.7

2.0

V

V

12.3

device's current draw is the sum of I_{CCWS} or I_{CCES} and I_{CCR} or I_{CCW}, respectively.

3. Includes RY/BY#.

Configuration Operations

Configuration Operations

V_{CC} Lockout Voltage

V_{CCW} during Block Erase, Full Chip

Erase, Word/Byte Write or Lock-Bit

4. Block erases, full chip erase, word/byte writes and lock-bit configurations are inhibited when V_{CCW}≤V_{CCWLK}, and not guaranteed in the range between V_{CCWLK}(max.) and V_{CCWH1}(min.), between V_{CCWH1}(max.) and V_{CCWH2}(min.) and above V_{CCWH2}(max.).

5. The Automatic Power Savings (APS) feature is placed automatically power save mode that addresses not switching more than 300ns while read mode.

6. About all of pin except describe Test Conditions, CMOS level inputs are either V_{CC}±0.2V or GND±0.2V, TTL level inputs are either V_{IL} or V_{IH} . 7. Sampled, not 100% tested.

Applying 12V±0.3V to V_{CCW} during erase/write can only be done for a maximum of 1000 cycles on each block. V_{CCW} may be connected to 12V±0.3V for a total of 80 hours maximum.

6.2.4 AC CHARACTERISTICS - READ-ONLY OPERATIONS⁽¹⁾

	$V_{CC}=2.7V-3.6V, T_{A}=-40^{\circ}C$ to	+85°C			
Sym.	Parameter	Notes	Min.	Max.	Unit
t _{AVAV}	Read Cycle Time		90		ns
t _{AVOV}	Address to Output Delay			90	ns
t _{ELOV}	CE# to Output Delay	2		90	ns
t _{PHOV}	RP# High to Output Delay			600	ns
t _{GLOV}	OE# to Output Delay	2		40	ns
t _{ELOX}	CE# to Output in Low Z	3	0		ns
t _{EHOZ}	CE# High to Output in High Z	3		40	ns
^t GLOX	OE# to Output in Low Z	3	0		ns
t _{GHOZ}	OE# High to Output in High Z	3		15	ns
t _{OH}	Output Hold from Address, CE# or OE# Change, Whichever Occurs First	3	0		ns
t _{FVOV}	BYTE# to Output Delay	3		90	ns
t _{FLOZ}	BYTE# Low to Output in High Z	3		25	ns
t _{ELFV}	CE# to BYTE# High or Low	3,4		5	ns

NOTES:

See AC Input/Output Reference Waveform for maximum allowable input slew rate.
 OE# may be delayed up to t_{ELQV}-t_{GLQV} after the falling edge of CE# without impact on t_{ELQV}.
 Sampled, not 100% tested.

4. If BYTE# transfer during reading cycle, exist the regulations separately.



LHF16J04



LHF16J04

6.2.5 AC CHARACTERISTICS - WRITE OPERATIONS⁽¹⁾

Sym.	Parameter	Notes	Min.	Max.	Unit
t _{AVAV}	Write Cycle Time		90		ns
t _{PHWL}	RP# High Recovery to WE# Going Low	2	1		μs
t _{elwl}	CE# Setup to WE# Going Low		10		ns
twlwh	WE# Pulse Width		50		ns
t _{SHWH}	WP#V _{IH} Setup to WE# Going High	2	100		ns
t _{VPWH}	V _{CCW} Setup to WE# Going High	2	100		ns
t _{AVWH}	Address Setup to WE# Going High	3	50		ns
t _{DVWH}	Data Setup to WE# Going High	3	50		ns
t _{WHDX}	Data Hold from WE# High		0		ns
t _{WHAX}	Address Hold from WE# High		0		ns
t _{WHEH}	CE# Hold from WE# High		10		ns
t _{WHWL}	WE# Pulse Width High		30		ns
t _{WHRL}	WE# High to RY/BY# Going Low or SR.7 Going "0"			100	ns
twhGL	Write Recovery before Read		0		ns
t _{QVVL}	V _{CCW} Hold from Valid SRD, RY/BY# High Z	2,4	0		ns
tovsl	WP# V _{IH} Hold from Valid SRD, RY/BY# High Z	2.4	0		ns
FVWH	BYTE# Setup to WE# Going High	5	50		ns
t _{WHFV}	BYTE# Hold from WE# High	5	90		ns

NOTES:

1. Read timing characteristics during block erase, full chip erase, word/byte write and lock-bit configuration operations are the same as during read-only operations. Refer to AC Characteristics for read-only operations.

2. Sampled, not 100% tested.

Complete, not 100/0 tested.
 Refer to Table 4 for valid A_{IN} and D_{IN} for block erase, full chip erase, word/byte write or lock-bit configuration.
 V_{CCW} should be held at V_{CCWH1/2} until determination of block erase, full chip erase, word/byte write or lock-bit configuration success (SR.1/3/4/5=0).
 IS DYTET

5. If BYTE# switch during reading cycle, exist the regulations separately.

LHF16J04



6.2.6 ALTERNATIVE CE#-CONTROLLED WRITES⁽¹⁾

Sym.	Parameter	Notes	Min.	Max.	Unit
t _{AVAV}	Write Cycle Time		90		ns
PHEL	RP# High Recovery to CE# Going Low	2	1		μs
WLEL	WE# Setup to CE# Going Low		0		ns
ELEH	CE# Pulse Width		70		ns
SHEH	WP#V _{IH} Setup to CE# Going High	2	100		ns
VPEH	V _{CCW} Setup to CE# Going High	2	100		ns
AVEH	Address Setup to CE# Going High	3	50		ns
DVEH	Data Setup to CE# Going High	3	50		ns
EHDX	Data Hold from CE# High		0		ns
EHAX	Address Hold from CE# High		0		ns
EHWH	WE# Hold from CE# High		0		ns
EHEL	CE# Pulse Width High		25		ns
EHRL	CE# High to RY/BY# Going Low or SR.7 Going "0"			100	ns
EHGL	Write Recovery before Read		0	L	ns
OVVL	V _{CCW} Hold from Valid SRD, RY/BY# High Z	2,4	0		ns
OVSL	WP# V _{IH} Hold from Valid SRD. RY/BY# High Z	2,4	0		ns
FVEH	BYTE# Setup to CE# Going High	5	50		ns
EHFV	BYTE# Hold from CE# High	5	90		ns

1. In systems where CE# defines the write pulse width (within a longer WE# timing waveform), all setup, hold, and inactive WE# times should be measured relative to the CE# waveform.

2. Sampled, not 100% tested.

Refer to Table 4 for valid A_{IN} and D_{IN} for block erase, full chip erase, word/byte write or lock-bit configuration.
 V_{CCW} should be held at V_{CCWH1/2} until determination of block erase, full chip erase, word/byte write or lock-bit configuration success (SR.1/3/4/5=0).

5. If BYTE# switch during reading cycle, exist the regulations separately.

LHF16J04



LHF16J04



NOTES:

1. If RP# is asserted while a block erase, full chip erase, word/byte write or lock-bit configuration operation is not executing, the reset will complete within 100ns.

2. A reset time, t_{PHQV}, is required from the later of RY/BY#(SR.7) going High Z("1") or RP# going high until outputs are

valid. Refer to AC Characteristics - Read-Only Operations for t_{PHQV} . 3. When the device power-up, holding RP# low minimum 100ns is required after V_{CC} has been in predefined range and also has been in stable there.

6.2.8 BLOCK ERASE, FULL CHIP ERASE, WORD/BYTE WRITE AND LOCK-BIT CONFIGURATION PERFORMANCE⁽³⁾

	T	<u> </u>		$V_{CCW} = 2.7V - 3.6V$		V _{CCW} =11.7V-12.3V		
Sym.	Param	eter	Notes	Typ. ⁽¹⁾	Max.	Typ. ⁽¹⁾	Max.	Unit
twhQV1	Word Write Time	32K word Block	2	33	200	20		μs
tEHQVI		4K word Block	2	36	200	27		μs
EnQVI	Byte Write Time	64K byte Block	2	31	200	19		μs
		8K byte Block	2	32	200	26		μs
	Block Write Time	32K word Block	2	1.1	4	0.66		S
	(In word mode)	4K word Block	2	0.15	0.5	0.12		S
	Block Write Time	64K byte Block	2	2.2	7	1.4		S
	(In byte mode)	8K byte Block	2	0.3	1	0.25		S
t _{WHQV2} t _{EHQV2}	Block Erase Time	32K word Block 64K byte Block	2	1.2	6	0.9		s
2		4K word Block 8K byte Block	2	0.6	5	0.5		S
	Full Chip Erase Time		2	42	210	32		s
t _{WHQV3} t _{EHQV3}	Set Lock-Bit Time Clear Block Lock-Bits Time		2	56	200	42		μs
t _{WHQV4} t _{EHOV4}			2	1	5	0.69		s
t _{WHRZ1}	Word/Byte Write Suspend Latency Time to Read		4	6	15	6	15	μs
t _{WHRZ2}	Block Erase Suspend Latency Time to Read		4	16	30	16	30	μs

 $V_{ac}=2.7V-3.6V$ T = -40°C to +85°C

NOTES:

Typical values measured at T_A=+25°C and V_{CC}=3.0V, V_{CCW}=3.0V or 12.0V. Assumes corresponding lock-bits are not set. Subject to change based on device characterization.

Excludes system-level overhead.
 Sampled but not 100% tested.

4. A latency time is required from issuing suspend command(WE# or CE# going high) until RY/BY# going High Z or SR.7 going "1".



.

	Refer to attach	ned drawing	
4. Stora	ge and Opening	of Dry Packing	
4-1.	Store under cor	nditions shown below before opening the dry packin	ng
	• •	ire range ∶ 5∼40℃	
	(2) Humidity	: 80% RH or less	
4-2.1	Notes on openir	ng the dry packing	
	(1) Before op	pening the dry packing, prepare a working table wh	nich is
	grounded	against ESD and use a grounding strap.	
	(2) The tray	has been treated to be conductive or anti-static.	. If the
	device is	s transferred to another tray, use a equivalent to	ray.
4-3.3	- Storage after o	opening the dry packing	
]	Perform the fol	lowing to prevent absorption of moisture after op	pening.
	(1) After ope	ening the dry packing, store the ICs in an environ	nment with a
	=	are of $5{\sim}25^\circ\!\mathrm{C}$ and a relative humidity of 60% or	less and
	mount ICs	within 72 hours after opening dry packing.	
4 - 4	Baking (drying)	before mounting	
	(1) Baking is	,	
	(A) If	the humidity indicator in the desiccant becomes p	pink
	(B) If	the procedure in section $4-3$ could not be perform	rmed
1	(2) Recommend	led baking conditions	
	If the ab	pove conditions (A) and (B) are applicable, bake i	it before
	mounting.	The recommended conditions are 16~24 hours at 12	20℃.
	Heat resi	stance tray is used for shipping tray.	
5 Surface	e Mount Conditi	ons	
		the following conditions when mounting ICs not to	o deteriorate IC
	quality.		
5 - 1.5	dering condit	ions (The following conditions are valid only for	one time solderin
	inting Method	Temperature and Duration	Measurement Poir
	low soldering	Peak temperature of 230℃ or less,	IC package
ł	(air)	duration of less than 15 seconds.	surface
		200° C or over, duration of less than 40 seconds.	
		Temperature increase rate of $1\sim4^{\circ}C/second$	
Mar	nual soldering	260℃ or less, duration of less	IC outer lead
(50	oldering iron)	than 10 seconds.	surface

5-2. Conditions for removal of residual flux

•	**									
	(1)	Ultrasonic	washing	power	:	25	Watts/liter	or	less	

- (1)Washing timeTotal 1 minute maximum(3)Solvent temperature $15 \sim 40^{\circ}$ C



mm



DRAWING NO.

AA1142

UNIT

mm





«Supplementary data»

LHF16J04					
Recommended mounting conditions for two time reflow soldering .					
Product name(Package)	LH28F160BJHE-TTL90(TSOP48-P-1220)				
Packing specification -	Tray (Dry packing)				
Mounting method	Reflow soldering (Air)				
Reflow soldering conditions	Peak temperature of 230℃ or less.				
	200°C or over, duration of less than 40 seconds.				
	Preheat temperature of 125∼150℃,duration of less				
	than 180 seconds. Temperature increase rate of				
	1~4°C/second.				
Measurement point	IC package surface				
Storage conditions	After opening the dry packing, store the ICs in				
	an environment with a temperature of $5\!\sim\!25^\circ\!\mathrm{C}$ and				
	a relative humidity of 60% or less.				
	If doing reflow soldering twice, do the first				
	reflow soldering within 72 hours after opening				
	dry packing and do the second reflow soldering				
	within 72 hours after the first reflow soldering.				
Note	If the above storage conditions are not				
	applicable, bake it before reflow soldering.				
	The recommended conditions are 16~24 hours				
	at 120°C.				
	(Heat resistance tray is used for shipping tray.)				

Recommended Reflow Soldering(Air) Temperature Profile



(NO. 990928-X15)

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